# Approximating the Visible Region of a Point on a Terrain* 

Boaz Ben-Moshe ${ }^{\dagger} \quad$ Paz Carmi ${ }^{\ddagger} \quad$ Matthew J. Katz ${ }^{\S}$


#### Abstract

Given a terrain $T$ and a point $p$ on or above it, we wish to compute the region $R_{p}$ that is visible from $p$. We present a generic radar-like algorithm for computing an approximation of $R_{p}$. The algorithm extrapolates the visible region between two consecutive rays (emanating from $p$ ) whenever the rays are close enough; that is, whenever the difference between the sets of visible segments along the cross sections in the directions specified by the rays is below some threshold. Thus the density of the sampling by rays is sensitive to the shape of the visible region. We suggest a specific way to measure the resemblance (difference) and to extrapolate the visible region between two consecutive rays. We also present an alternative algorithm, which uses circles of increasing radii centered at $p$ instead of rays emanating from $p$. Both algorithms compute a representation of the (approximated) visible region that is especially suitable for visibility from $p$ queries. Finally, we report on the experiments that we performed with these algorithms and with their corresponding fixed versions, using a natural error measure. Our main conclusion is that the radar-like algorithm is significantly better than the others.


## 1 Introduction

Let $T$ be a triangulation representing a terrain (i.e., there is a height ( $z$-coordinate) associated with each triangle vertex). We are interested in the following well known problem. Given a point $p$ on (or above) $T$, compute the region $R_{p}$ of $T$ that is visible from $p$. A point $q$ on $T$ is visible from $p$ if and only if the line segment $\overline{p q}$ lies above $T$ (in the weak sense). Thus $R_{p}$

[^0]consists of all points on $T$ that are visible from $p$.
The problem of computing the visible region of a point arises as a subproblem in numerous applications (see, e.g., $[3,7,9,11]$ ), and, as such, has been studied extensively $[2,3,4,5,7]$. For example, the coverage area of an antenna for which line of sight is required may be approximated by clipping the region that is visible from the tip of the antenna with an appropriate disk centered at the antenna.

Since the combinatorial complexity of $R_{p}$ might be $\Omega\left(n^{2}\right)[3,6]$, where $n$ is the number of triangles in $T$, it is desirable to also have fast approximation algorithms, i.e., algorithms that compute an approximation of $R_{p}$. Moreover, a good approximation of the visible region is often sufficient, especially when the triangulation itself is only a rough approximation of the underlying terrain. Note that in this paper we are assuming that the terrain representation (i.e., the triangulation $T$ ) is fixed and cannot be modified. Simplifying the triangulation can of course result in a significant decrease in the actual running time of any algorithm for computing the visible region. This approach was studied in a previous paper [1]. See, e.g., [8] for more information on terrain simplification.

We present a generic radar-like algorithm for computing an approximation of $R_{p}$. The algorithm computes the visible segments along two rays $\rho_{1}, \rho_{2}$ emanating from $p$, where the angle between the rays is not too big. (I.e., each of the rays specifies a direction, and the algorithm computes the (projections of the) visible portions of the cross section of $T$ in this direction.) It then has to decide whether the two sets of visible segments (one per ray) are close enough so that it can extrapolate the visible region of $p$ within the wedge defined by $\rho_{1}$ and $\rho_{2}$, or whether an intermediate ray is needed. In the latter case the algorithm will now consider the smaller wedge defined by $\rho_{1}$ and the intermediate ray. Thus a nice property of the algorithm is that the density of the sample rays varies and depends on the shape of $R_{p}$.

In order to use this generic algorithm one must provide (i) a measure of resemblance for two sets of visible segments, where each set consists of the visible segments along some ray from $p$, and (ii) an algorithm to extrapolate the visible region between two rays whose corresponding sets were found similar enough.

In Section 2 we describe in more detail the generic algorithm and provide the missing ingredients.

In Section 3 we present several other algorithms for computing the visible region $R_{p}$. The first algorithm computes $R_{p}$ exactly. Since we need such an algorithm for the experimental evaluation of our approximate algorithms, we decided to devise one that is based on the general structure of the radar-like algorithm. Our exact algorithm is rather simple and is based on known results; nevertheless it seems useful. Specifically, the algorithm repeatedly computes the portion of $R_{p}$ within a slice, defined by a pair of rays passing through vertices of the terrain, that does not contain a vertex of $T$ in its interior. This computation can be done efficiently as is shown in [5].

The second algorithm (called the expanding circular-horizon algorithm or ECH for short) is in some sense orthogonal to the radar-like algorithm; it uses circles of increasing radii centered at the view point $p$ instead of rays emanating from $p$. It is influenced by the exact algorithm described by De Floriani and P. Magillo [5]. The algorithm approximates the visible region $R_{p}$ by maintaining the (approximate) viewing angles corresponding to a set of sample points on the expanding circular front (see Section 3). This allows us to partition the current front into maximal visible and invisible arcs. We now examine the sets of visible arcs on the current and previous fronts. If they are close enough, then the portion of $R_{p}$ within the annulus defined by the two circles is approximated. Otherwise, we compute the visible arcs on a circle of intermediate radius and repeat.

Both the radar-like algorithm and the expanding circular-horizon algorithm have corresponding fixed versions, that play an important role in the experiments that were performed, see below. In the fixed version of the radar-like algorithm, the angle between two consecutive rays is fixed and we approximate the portion of $R_{p}$ in the sector defined by the rays in any case, even if they are not close enough. In the fixed version of the expanding circular-horizon algorithm, the increase in radius between two consecutive circles is fixed and again we approximate the portion of $R_{p}$ in the annulus defined by the circles in any case.

In Section 4 we suggest a natural way to measure the error in an approximation $R_{p}^{\prime}$ of $R_{p}$ produced by one of our algorithms. The error associated with $R_{p}^{\prime}$ is the area of the XOR of $R_{p}^{\prime}$ and $R_{p}$, divided by the area of the disk of radius $l$, where $l$ is the range of sight that is in use. Using this error measure (and the exact algorithm), we performed a collection of experiments (described in Section 4) with the radar-like and expanding circular-horizon algorithms and their corresponding
fixed versions. Our main conclusions from these experiments are that (i) the sensitive versions are significantly better than their corresponding fixed versions (when the total number of slices / annuli composing the final approximation is the same in both versions), and (ii) the radar-like algorithm is significantly better than the expanding circular-horizon algorithm. In Section 4 we offer some explanations to these findings.

## 2 The Radar-Like Algorithm

In this section we first present our radar-like generic algorithm. Next we describe the measure of resemblance and the extrapolation algorithm that we devised, and that are needed in order to transform the generic algorithm into an actual algorithm.

The generic algorithm is presented in the frame below. The basic operation that is used is the crosssection operation, denoted cross-section $(T, p, \theta)$, which computes the visible segments along the ray emanating from $p$ and forming an angle $\theta$ with the positive $x$ axis. More precisely, cross-section $(T, p, \theta)$ computes the (projections of the) visible portions of the cross section of $T$ in the direction specified by this ray. Roughly speaking, the generic algorithm sweeps the terrain $T$ counter clockwise with a ray $\rho$ emanating from $p$, performing the cross-section operation whenever the pattern of visible segments on $\rho$ is about to change significantly with respect to the pattern that was found by the previous call to cross-section. The algorithm then extrapolates, for each pair of consecutive patterns, the visible region of $p$ within the wedge defined by the corresponding locations of $\rho$.

Given a triangulation $T$ representing a terrain (i.e., with heights associated with the triangle vertices), and a view point $p$ on or above $T$ :
$\theta \leftarrow 0$.
$\alpha \leftarrow$ some constant angle, say, $\pi / 45$.
$S_{1} \leftarrow \operatorname{cross-section}(T, p, \theta)$.
$S_{2} \leftarrow \operatorname{cross-section}(T, p, \theta+\alpha)$.
while $(\theta \leq 360)$
if ( $S_{1}$ is close-enough to $S_{2}$ )
extrapolate $\left(S_{1}, S_{2}\right)$;
$\theta \leftarrow S_{2}$.angle;
$S_{1} \leftarrow S_{2}$;
$S_{2} \leftarrow \operatorname{cross}-\operatorname{section}(T, p, \min (\theta+\alpha, 360)) ;$ else
$\mu \leftarrow\left(S_{1}\right.$.angle $+S_{2}$.angle $) / 2 ;$
$S_{2} \leftarrow \operatorname{cross-section}(T, p, \mu)$;


Figure 1: Grey marks visible and black marks invisible. (a) The close-enough threshold function: $\delta$ times the relative length of the XOR of $S_{1}$ and $S_{2}$. (b) The extrapolate function.

In order to obtain an actual algorithm we must provide precise definitions of close-enough and extrapolate.
Close-enough: A threshold function that checks whether two patterns $S_{1}, S_{2}$ are similar, where each of the patterns corresponds to the set of visible segments on some ray from $p$. There are of course many ways to define close-enough. We chose the following definition. In practice, the rotating ray is actually a rotating segment of an appropriate length. Let $l$ denote this length. We refer to $l$ as the range of sight. Now rotate the ray containing $S_{2}$ clockwise until it coincides with the ray containing $S_{1}$. See Figure 1 (a). Next compute the length of the XOR of $S_{1}$ and $S_{2}$, that is, the total length covered by only one of the sets $S_{1}, S_{2}$. This length is then divided by $l$. Denote by $v$ the value that was computed, and let $\delta$ be the angle between $S_{1}$ and $S_{2}$. If $\delta \cdot v \leq C$, where $C$ is some constant, then return TRUE else return FALSE. The role of $\delta$ in the above formula is to force close-enough to return TRUE when the angle between the rays is small, even if the patterns that are being compared still differ significantly.
Extrapolate: Given two patterns $S_{1}, S_{2}$ which are close-enough, we need to compute an approximation of the portion of the visible region of $p$ that is contained in the corresponding wedge. We do this as follows. Consider Figure 1 (b). For each 'event point' (i.e., start or end point of a visible segment) on one of the two horizontal rays, draw a vertical segment that connects it with the corresponding point on the other ray. For each rectangle that is obtained color it as follows, where grey means visible and black means invisible. If the horizontal edges of a rectangle are either both visible
from $p$ or both invisible from $p$, then, if both are visible, color it grey, else color it black. If, however, one of the horizontal edges is visible and the other is invisible, divide the rectangle into four pieces by drawing the two diagonals. The color of the upper and lower pieces is determined by the color of the upper and lower edges, respectively, and the color of the left and right pieces is determined by the color of the rectangles on the left and on the right, respectively. Assuming there are no two event points such that one of them is exactly above the other, the coloring procedure is well defined. That is, the odd numbered rectangles will be colored with a single color, and the even numbered rectangles will be divided.
Remark 1. The representation of the (approximated) visible region $R_{p}^{\prime}$ that is computed by the radar-like algorithm is especially suitable for queries of the form: Given a query point $q$ on $T$, determine whether $q$ is visible from $p$, or, more precisely, determine whether (the projection of) $q$ lies in $R_{P}^{\prime}$. We first verify that $q$ is within the range of sight $l$, i.e., that $q$ lies within the disk of radius $l$ centered at $p$. Next we determine whether $q$ lies in $R_{p}^{\prime}$, in logarithmic time, by two binary searches. The first search locates the sector of the disk in which $q$ lies, and the second locates the 'rectangle' within the sector in which $q$ lies. Finally, it remains to check in which of the at most four triangles corresponding to this rectangle $q$ lies.
Remark 2. One can think of alternative definitions for close-enough and extrapolate. However, it seems reasonable to require the following two properties: (i) A small change in the set of visible segments along a ray should only cause small changes in the close-
enough measure and in the visible region computed by extrapolate (within the appropriate wedge), and (ii) If there are no "surprises" between two close enough rays, then the visible region computed by extrapolate within the wedge should be very similar to the real visible region. In addition, the definitions should remain simple and easy to implement.

## 3 Other Algorithms

In this section we present several other algorithms for computing the visible region. The first algorithm computes the visible region exactly; its general structure is similar to that of the radar-like algorithm. The second algorithm (called the expanding circular-horizon algorithm or ECH for short) is influenced by the exact algorithm of De Floriani and Magillo [5]. It computes an approximation of the visible region using circles of increasing radii (instead of rays) and similar definitions of close-enough and extrapolate. Towards the end of this section we mention the fixed versions of the radar-like and expanding circular-horizon algorithms.

The algorithms presented in this section are part of our testing environment for the radar-like algorithm. However, we believe that the exact algorithm and the expanding circular-horizon algorithm are of independent interest.
3.1 The exact algorithm. Since we need an exact algorithm (i.e., an algorithm that computes the visible region $R_{p}$ exactly) for the experimental evaluation of our approximate algorithms, we decided to devise one that is based on the general structure of the radar-like algorithm, instead of using one of the known algorithms. Our exact algorithm is rather simple. It repeatedly computes the (exact) portion of $R_{p}$ within a slice that is defined by a pair of rays passing through vertices of the terrain $T$, and that does not contain a vertex of $T$ in its interior; see Figure 2. This can be done in time $O(m \log m)$, where $m$ is the number of edges of the terrain that cross the slice, as is shown in [5].
Remark. The radar-like algorithm can be modified accordingly, so that whenever the slice under consideration does not contain a vertex of $T$ in its interior, the portion of $R_{p}$ within the slice is computed exactly, rather than calling close-enough (possibly more than once) and then extrapolate.

### 3.2 The expanding circular-horizon algorithm.

This algorithm is in some sense orthogonal to the radarlike algorithm; it uses circles of increasing radii centered at the view point $p$ instead of rays emanating from $p$.

For a point $q$ on $T$, let $\theta_{q}$ be the minimum viewing angle (from $p$ ) at which a vertical pole of infinite
height based at $q$ can be seen from $p ; \theta_{q}$ is the viewing angle corresponding to $q$. Our algorithm approximates the visible region $R_{p}$, by maintaining the (approximate) viewing angles corresponding to the points on the expanding circular horizon. More precisely, the algorithm only considers the points on the circular horizon at directions $\alpha, 2 \alpha, 3 \alpha, \ldots$ with respect to $p$, where $\alpha$ is a parameter of the algorithm. Initially, the circular horizon is the point $p$ itself, and the corresponding viewing angle is $-\pi / 2$.


Figure 3: Left: $\theta_{q}=\phi_{q}$ and $q$ is visible; Right: $\theta_{q}=\theta_{q^{\prime}}$ and $q$ is invisible.

The (approximate) viewing angles for the current circular horizon are computed from those of the previous circular horizon (by resolve-viewing-angles) as follows (see Figure 3). Let $q$ be a point on the current horizon at direction $i \alpha$ with respect to $p$, and let $q^{\prime}$ be the point on the previous horizon at the same direction. Let $\phi_{q}$ be the angle between $p q$ and the horizontal plane through $p$, then $\theta_{q} \leftarrow \max \left\{\theta_{q^{\prime}}, \phi_{q}\right\}$, and $q$ is said to be visible from $p$ if and only if $\theta_{q}=\phi_{q}$.

After applying resolve-viewing-angles to the current horizon $C$, we partition $C$ into maximal visible and invisible arcs as follows. An arc of $C$ between two consecutive sample points (i.e., points at direction $i \alpha$ and $(i+1) \alpha$ with respect to $p$ ) is called a primitive arc. We first consider each of the primitive arcs $a$. If both endpoints of $a$ are visible (resp., invisible), then we assume all points in $a$ are visible (resp., invisible). If however one of the endpoints $q_{i}$ is visible and the other $q_{i+1}$ is invisible, we assume all points in the half of $a$ adjacent to $q_{i}$ are visible and all points in the half adjacent to $q_{i+1}$ are invisible. We now can partition $C$ into maximal visible and invisible arcs.

We consider the current and previous horizons $C_{2}$ (of radius $r_{2}$ ) and $C_{1}$ of radius ( $r_{1}$ ), respectively. As in the radar-like algorithm, we must decide whether the two sets of visible arcs are close-enough. If yes, we call extrapolate to approximate the portion of $R_{p}$ that is contained in the annulus defined by $C_{1}$ and $C_{2}$. If no, we generate the intermediate horizon $C$ of radius $\left(r_{1}+r_{2}\right) / 2$, and repeat for the pair $C_{1}, C$. We use similar definitions of close-enough and extrapolate as in the radar-like algorithm.


Figure 2: Left: the exact algorithm draws a ray through each vertex of $T$; Right: a slice that is defined by two consecutive rays, the corresponding cross sections, and the exact portion of $R_{p}$ within the slice.

Given a triangulation $T$ representing a terrain (i.e., with heights associated with the triangle vertices), and a view point $p$ on or above $T$ :
$\alpha \leftarrow$ some constant angle, say $\pi / 180$.
$d \leftarrow$ some constant distance, say, 10 meters.
$r_{1} \leftarrow r_{\text {min }}$.
$C_{1} \leftarrow$ determine the viewing angles corresponding to the $2 \pi / \alpha$ sample points on the circle of radius $r_{1}$. $C_{2} \leftarrow$ resolve-viewing-angles $\left(T, p, C_{1}, r_{1}+d\right)$. while ( $r_{1} \leq r_{\text {max }}$ )
if ( $C_{1}$ is close-enough to $C_{2}$ )
extrapolate $\left(C_{1}, C_{2}\right)$;
$r_{1} \leftarrow C_{2}$.radius;
$C_{1} \leftarrow C_{2}$;
$r \leftarrow \min \left(r_{1}+d, r_{\text {max }}\right) ;$
$C_{2} \leftarrow$ resolve-viewing-angles $\left(T, p, C_{1}, r\right)$;
else
$r \leftarrow\left(C_{1}\right.$.radius $+C_{2}$.radius $) / 2 ;$
$C_{2} \leftarrow$ resolve-viewing-angles $\left(T, p, C_{1}, r\right)$;
3.3 The corresponding fixed versions. Both the radar-like algorithm and the expanding circular-horizon algorithm have corresponding fixed versions. In the fixed version of the radar-like algorithm the angle between two consecutive rays is fixed and we approximate the portion of $R_{p}$ in the sector defined by the rays in any case, even if they are not close-enough; see Figure 4. In the fixed version of the expanding circular-horizon algorithm, the increase in radius between two consecutive
circles is fixed and again we approximate the portion of $R_{p}$ in the annulus defined by the circles in any case, see Figure 5.

## 4 Experimental Results

In this section we report on the experiments that we performed with the approximation algorithms described in Sections 2 and 3. Namely, the radar-like algorithm, the expanding circular-horizon algorithm (ECH), and their corresponding fixed versions. We have also implemented the exact algorithm (Section 3.1), which is needed for the error computation.
4.1 The error measure. In our experiments we use the following natural error measure. Let $R_{p}^{\prime}$ be an approximation of $R_{p}$ obtained by some approximation algorithm, where $R_{p}$ is the region visible from $p$. Then the error associated with $R_{p}^{\prime}$ is the area of the XOR of $R_{p}^{\prime}$ and $R_{p}$, divided by the area of the disk of radius $l$, where $l$ is the range of sight that is in use. See Figure 6.
4.2 The experiments. Ten input terrains representing ten different and varied geographic regions were used. Each input terrain covers a rectangular area of approximately $15 \times 10 \mathrm{~km}^{2}$, and consists of approximately $5,000-10,000$ triangle vertices. For each terrain we picked several view points ( $x, y$ coordinates) randomly. For each view point $p$ we applied each of the four approximation algorithms (as well as the exact algorithm) 20 times: once for each combination of height (either $1,10,20$, or 50 meters above the surface of $T$ ) and range of sight (either 500, 1000, 1500, 2500, or 3500 meters). For each (approximated) region that was obtained, we computed the associated error, according to


Figure 4: The visible region computed by the radar-like algorithm (left) and by its corresponding fixed version (right), each composed of 72 slices.


Figure 5: The visible region computed by the expanding circular-horizon algorithm (left) and by its corresponding fixed version (right).


Figure 6: Left: the exact region $R_{p}$; Middle: the approximate region $R_{p}^{\prime}$ computed by the radar-like algorithm; Right: $\operatorname{XOR}\left(R_{p}^{\prime}, R_{p}\right)$.
the error measure above. All this was repeated three times; once per each of three levels of sampling (see below).

The level of sampling is determined by the number of calls to extrapolate that are issued during the execution of an algorithm. We used three levels of sampling: 80, 140, and 220 . Since the extrapolation between two consecutive rays is comparable to the extrapolation between two circular horizons, this seems a fair basis for comparison. (In order to achieve a specific level of sampling when running one of the nonfixed versions, we repeated the computation several times, with different values of the constant $C$, until the desired level of sampling was reached.)

| Accuracy level: 80 | $\mathrm{l}=500$ | $\mathrm{l}=1000$ | $\mathrm{l}=1500$ | $\mathrm{l}=2500$ | $\mathrm{l}=3500$ |
| :--- | :---: | :---: | :---: | :---: | :---: |
| Fixed ECH | 4.53 | 3.97 | 3.87 | 3.88 | 3.75 |
| ECH | 3.71 | 3.49 | 3.40 | 3.29 | 3.35 |
| Fixed radar-like | 2.09 | 1.36 | 1.19 | 1.08 | 0.97 |
| Radar-like | 1.37 | 0.91 | 0.78 | 0.65 | 0.62 |

Table 1: Results for sampling level 80.

| Accuracy level: 140 | 500 | 1000 | 1500 | 2500 | 3500 |
| :--- | :---: | :---: | :---: | :---: | :---: |
| Fixed ECH | 2.71 | 2.35 | 2.12 | 2.16 | 2.12 |
| ECH | 2.30 | 2.09 | 1.88 | 1.88 | 1.94 |
| Fixed radar-like | 1.21 | 0.97 | 0.91 | 0.81 | 0.72 |
| Radar-like | 0.89 | 0.71 | 0.62 | 0.59 | 0.53 |

Table 2: Results for sampling level 140.
4.3 The results. Our results are presented in the following two sets of tables. The first three tables show the error for each of the four algorithms as a function of the sampling level and the range of sight. Consider,

| Accuracy level: 220 | 500 | 1000 | 1500 | 2500 | 3500 |
| :--- | :---: | :---: | :---: | :---: | :---: |
| Fixed ECH | 1.36 | 1.30 | 1.18 | 1.19 | 1.20 |
| ECH | 1.22 | 1.15 | 1.13 | 1.09 | 1.13 |
| Fixed radar-like | 0.79 | 0.63 | 0.59 | 0.51 | 0.40 |
| Radar-like | 0.53 | 0.41 | 0.37 | 0.29 | 0.28 |

Table 3: Results for sampling level 220.
for example, the first table. This table contains our results for sampling level 80 . The first line in this table corresponds to the fixed version of the expanding circular-horizon algorithm (ECH). The first entry in this line (4.53) is the average error (in percents) obtained when running ECH with accuracy level 80 and range of sight 500 for each of the view points (over all terrains) and each of the four possible heights.

Tables 4 and 5 show the amount of work needed in order to reach a certain level of accuracy. In Table 4 the amount of work is measured by the number of calls to cross-section (alternatively, resolve-viewing-angles), and in Table 5 the amount of work is measured by the total running time. For example, using the fixed radar-like algorithm, the average number of calls to cross-section needed to obtain an error of 1 percent was 80 , and, using the ECH algorithm, the average running time needed to obtain an error of 0.5 percent was 1648 milliseconds. All experiments were performed on the following platform: Pentium $4,2.4 \mathrm{GHz}$, 512MB, Linux 8.1, Java 1.4.
4.4 Conclusions. Based on the results above the radar-like approach is significantly better than the expanding circular-horizon approach. For each of the sampling levels, the regions computed by the two radar-like algorithms were more accurate than those computed by the two ECH algorithms for any range of sight (see Ta-

| Error: | 1.00 | 0.75 | 0.50 |
| :--- | :---: | :---: | :---: |
| Fixed ECH | 263 | 616 | 1009 |
| ECH | 231 | 522 | 893 |
| Fixed radar-like | 80 | 140 | 220 |
| Radar-like | 61 | 103 | 174 |

Table 4: Average number of calls to cross-section (alternatively, resolve-viewing-angles) by accuracy level.

| Error: | 1.00 | 0.75 | 0.50 |
| :--- | :---: | :---: | :---: |
| Fixed ECH | 597 | 1045 | 1648 |
| ECH | 579 | 1012 | 1591 |
| Fixed radar-like | 112 | 192 | 301 |
| Radar-like | 101 | 168 | 274 |

Table 5: Average running time (in milliseconds) by accuracy level.
bles 1-3). Moreover, for each level of accuracy, the ECH algorithms had to work much harder than the radarlike algorithms (according to both measures) in order to reach the desired level of accuracy (see Tables 4-5).

A possible explanation for the better performance of the radar-like algorithms is that in the ECH algorithms the computation of the visible arcs on the current circular horizon is only an approximation, while in the radar-like algorithms the visible segments on a ray are computed exactly. Referring to Figure 7, if the ECH algorithms miss a ridge like the one in the left picture (drawn as a narrow rectangle), then all subsequent circles will miss it and therefore might conclude that the corresponding arcs are visible while they are not. On the other hand, a ridge like the one in the right picture that is missed by the radar-like algorithms does not affect subsequent computations leading to smaller errors.

Another clear conclusion is that the adaptive (i.e., non-fixed) versions are more accurate than their corresponding fixed versions. For each sampling level, the regions computed by the adaptive versions were significantly more accurate than those of the fixed versions (see Tables 1-3). This advantage of the adaptive versions is especially noticeable when the sampling level is low.

As expected, the adaptive versions are somewhat slower than the corresponding fixed versions for a given sampling level, since the adaptive versions perform more cross-section (alternatively, resolve-viewing-angles) operations. Actually, we found that on average the radarlike algorithm issues about 9 percent more calls to crosssection than the fixed radar-like algorithm. However,
when taking into consideration the improved accuracy of the adaptive version, we see (Table 5) that the the adaptive version is on average about 10 percent faster than the corresponding fixed versions.


Figure 7: ECH vs. the radar-like algorithm.

Finally, we recently performed some experiments with the two radar-like algorithms using somewhat larger terrains consisting of approximately 40,000 vertices and covering a rectangular area of approximately $20 \times 20 \mathrm{~km}^{2}$. In general, the errors that we got were somewhat smaller (using the same levels of sampling and ranges of sight), and the adaptive version remained more accurate than the fixed version. The smaller errors are probably due to the higher resolution, although, in general, the nature of the terrain can significantly affect the accuracy of the approximations.

Acknowledgment. The authors wish to thank Ofir Ganani and Maor Mishkin who helped implementing the radar-like and the expanding circular-horizon algorithms, and Joe Mitchell for helpful discussions.

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[^0]:    *Research by Ben-Moshe and Katz is partially supported by grant no. 2000160 from the U.S.-Israel Binational Science Foundation, and by the MAGNET program of the Israel Ministry of Industry and Trade (LSRT consortium). Research by Carmi is partially supported by a Kreitman Foundation doctoral fellowship.
    ${ }^{\dagger}$ Department of Computer Science, Ben-Gurion University of the Negev, Beer-Sheva 84105, Israel, benmoshe@cs.bgu.ac.il.
    ${ }^{\ddagger}$ Department of Computer Science, Ben-Gurion University of the Negev, Beer-Sheva 84105, Israel, carmip@cs.bgu.ac.il.
    ${ }^{\S}$ Department of Computer Science, Ben-Gurion University of the Negev, Beer-Sheva 84105, Israel, carmip@cs.bgu.ac.il.

